

NOM :

PRENOM :

UE Sciences de l'Information 1 - EC Modélisation des SED

Durée : 1h (documents non autorisés)

Les réponses doivent être saisies sur le sujet.

Question 1. Donnez les automates reconnaissant les expressions régulières suivantes (LG : langage généré et LM : langage marqué) (3 points)

		Points
$LG = \varepsilon + (ab)^*(\varepsilon + a + ac^*)$ $LM = (ab)^*ac^*$		
$LG = \varepsilon + a + abc^* + bc^*$ $LM = abc^* + bc^*$		

Question 2. A l'aide du lemme d'Aden, déterminez quels sont les langages générés et marqués par l'automate suivant (4 points)

		Points
	L1 = L2 = L3 = Par substitution, on obtient L1 = Par application du lemme d'Arden, on obtient alors : L1 = L2 = L3 = d'où Langage généré $L_G =$ Langage marqué $L_M =$	

Question 3. Déterminer l'automate suivant (4 points)

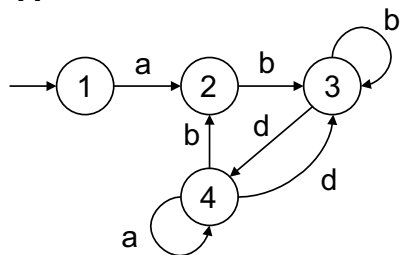
		Points

PRENOM :

Question 4. *Calculer la composition synchrone entre les automates A et B. (4 points)*

Points

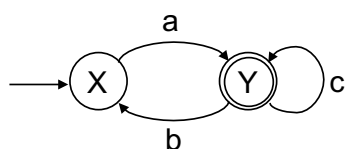
A



A X B

A II B

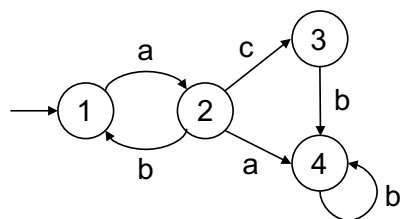
B



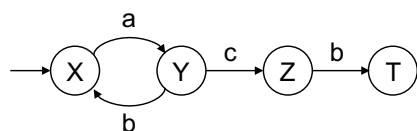
Question 5. Soit b un événement incontrôlable, et a et c deux événements contrôlables. Appliquez l'algorithme de Kumar pour déterminer le superviseur maximal permissif pour le procédé G et la spécification K . (5 points)

Points

G (Procédé)



K (spécification)



G X K

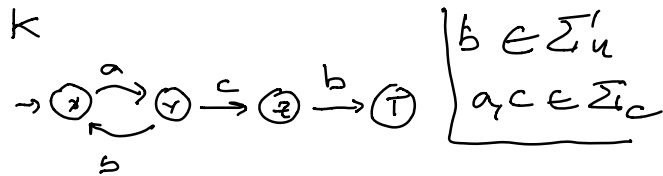
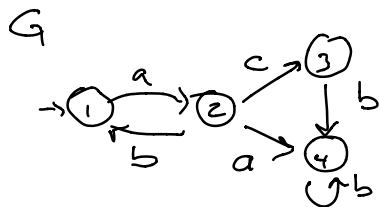
Etat(s) interdit(s) :

Justification :

Etat(s) faiblement interdit(s) :

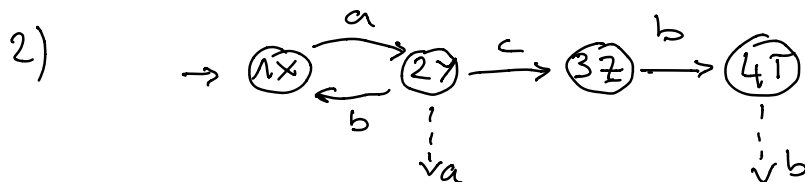
Justification :

Automate générant le langage suprême contrôlable $L(K/G)$



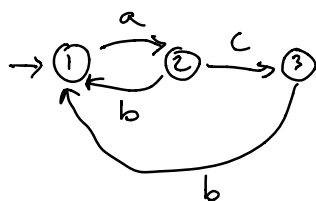
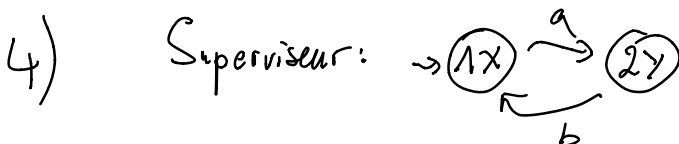
1) $\epsilon + (ab)^* (\epsilon + a + ac + acb^* + aab^*)$ $\epsilon + (ab)^* (\epsilon + a + ac + acb)$ $\underbrace{b}_{\Sigma'_u}$

$acbb \begin{cases} \in \overline{K} \Sigma'_u \cap L_G \\ \notin \overline{K} \end{cases}$



3) 4τ état interdit $\begin{cases} b \text{ admis ds } G \\ b \text{ interdit ds } K \\ b \in \Sigma'_u \end{cases}$

3ζ état forcément interdit $\begin{cases} 3\zeta \xrightarrow{b} 4\tau \\ b \in \Sigma'_u \end{cases}$



$$\begin{cases} L1 = \epsilon + L2b + L3b \\ L2 = L1a \\ L3 = L2c \end{cases}$$

(2)

$$\begin{aligned} L1 &= \epsilon + L1ab + L1acb \\ &= \epsilon + L1(ab + acb) \\ L1 &= (ab + acb)^* \\ L2 &= (ab + acb)^* a \end{aligned}$$

$$L3 = (ab + acb)^* ac$$

$$\begin{cases} L1 = \epsilon + L3a \end{cases}$$

$$\begin{cases} L2 = L1c \end{cases}$$

$$\begin{cases} L3 = L2a + L3b \end{cases}$$

(2 bis)

$$L3 = (\epsilon + L3a)ca + L3b$$

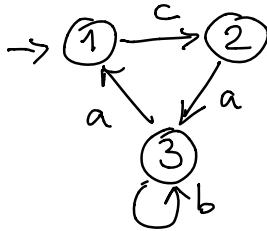
$$= ca + L3aca + L3b$$

$$= ca + L3(aca + b)$$

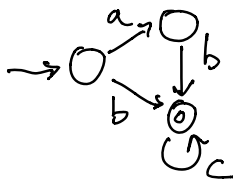
$$L3 = ca(aca + b)^*$$

$$L1 = \epsilon + ca(aca + b)^* a$$

$$L2 = c + ca(aca + b)^* ac$$



$$L = \epsilon + c + ca(aca + b)^*(\epsilon + a + ac)$$

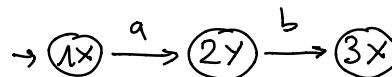
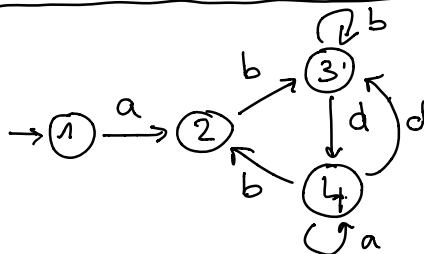


$$L_g = \epsilon + a + abc^* + bc^*$$

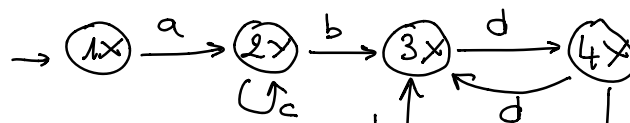
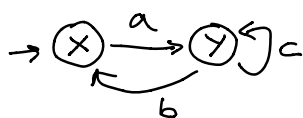
$$L_n = abc^* + bc^*$$

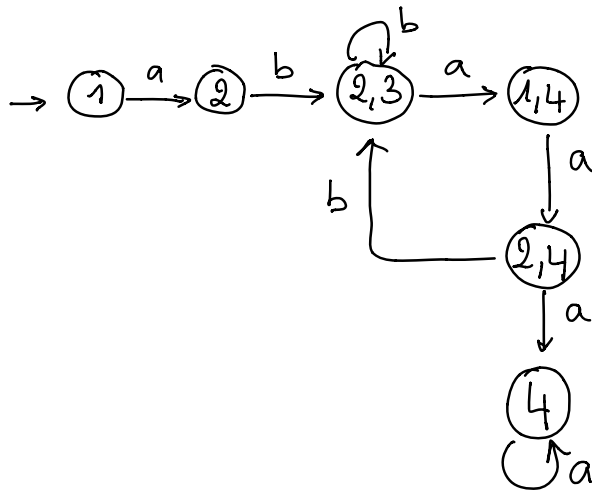
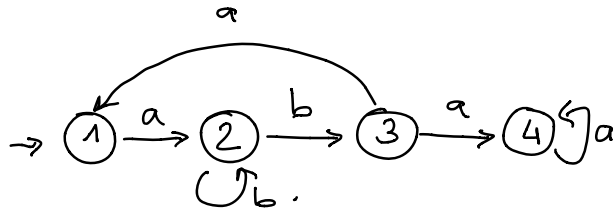
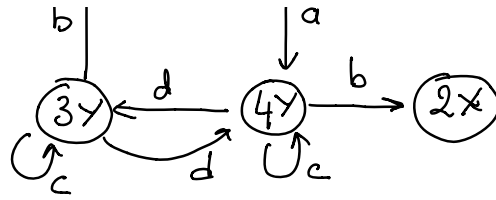
(1)

$$L_g = \epsilon + (ab)^*(\epsilon + a + ac^*)$$



(3)





$$\cdot 2 \xrightarrow{b} 2,3$$

$$\begin{cases} 2,3 \xrightarrow{b} 2,3 \\ 2,3 \xrightarrow{a} 1,4 \end{cases}$$

$$\cdot 1,4 \xrightarrow{a} 2,4$$

$$\begin{cases} 2,4 \xrightarrow{a} 4 \\ 2,4 \xrightarrow{b} 2,3 \end{cases}$$

$$\cdot 4 \xrightarrow{a} 4$$

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